



Chapter 6

Recursion as a Problem-Solving Technique

CS102 Sections 51 and 52

Marc Smith and Jim Ten Eyck
Spring 2008

© 2006 Pearson Addison-Wesley. All rights reserved

6-1

Backtracking

- Backtracking
 - A strategy for guessing at a solution and backing up when an impasse is reached
- Recursion and backtracking can be combined to solve problems

© 2006 Pearson Addison-Wesley. All rights reserved

6-2

The Eight Queens Problem

- Problem
 - Place eight queens on the chessboard so that no queen can attack any other queen
- Strategy: guess at a solution
 - There are 4,426,165,368 ways to arrange 8 queens on a chessboard of 64 squares

© 2006 Pearson Addison-Wesley. All rights reserved

6-3

The Eight Queens Problem

- An observation that eliminates many arrangements from consideration
 - No queen can reside in a row or a column that contains another queen
 - Now: only 40,320 arrangements of queens to be checked for attacks along diagonals

© 2006 Pearson Addison-Wesley. All rights reserved

6-4

The Eight Queens Problem

- Providing organization for the guessing strategy
 - Place queens one column at a time
 - If you reach an impasse, backtrack to the previous column

© 2006 Pearson Addison-Wesley. All rights reserved

6-5

The Eight Queens Problem

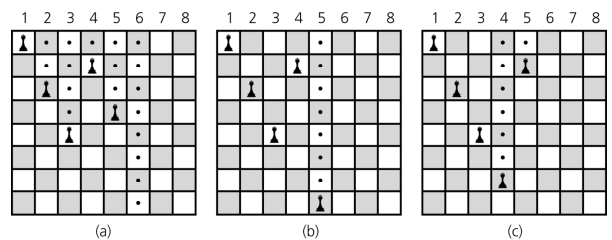


Figure 6-1

a) Five queens that cannot attack each other, but that can attack all of column 6; b) backtracking to column 5 to try another square for the queen; c) backtracking to column 4 to try another square for the queen and then considering column 5 again

© 2006 Pearson Addison-Wesley. All rights reserved

6-6

The Eight Queens Problem

- A recursive algorithm that places a queen in a column
 - Base case
 - If there are no more columns to consider
 - You are finished
 - Recursive step
 - If you successfully place a queen in the current column
 - Consider the next column
 - If you cannot place a queen in the current column
 - You need to backtrack

© 2006 Pearson Addison-Wesley. All rights reserved

6-7

The Eight Queens Problem

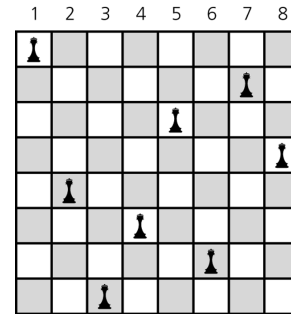


Figure 6-2
A solution to the Eight Queens problem

© 2006 Pearson Addison-Wesley. All rights reserved

6-8

Defining Languages

- A language
 - A set of strings of symbols
 - Examples: English, Java
 - If a Java program is one long string of characters, the language JavaPrograms is defined as
$$\text{JavaPrograms} = \{\text{strings } w : w \text{ is a syntactically correct Java program}\}$$

© 2006 Pearson Addison-Wesley. All rights reserved

6-9

Defining Languages

- A language does not have to be a programming or a communication language
 - Example
 - The set of algebraic expressions
$$\text{AlgebraicExpressions} = \{w : w \text{ is an algebraic expression}\}$$

© 2006 Pearson Addison-Wesley. All rights reserved

6-10

Defining Languages

- Grammar
 - States the rules for forming the strings in a language
- Benefit of recursive grammars
 - Ease of writing a recognition algorithm for the language
 - A recognition algorithm determines whether a given string is in the language

© 2006 Pearson Addison-Wesley. All rights reserved

6-11

The Basics of Grammars

- Symbols used in grammars
 - $x | y$ means x or y
 - $x y$ means x followed by y
(In $x \cdot y$, the symbol \cdot means concatenate, or append)
 - $\langle \text{word} \rangle$ means any instance of word that the definition defines

© 2006 Pearson Addison-Wesley. All rights reserved

6-12

The Basics of Grammars

- Java identifiers
 - A Java identifier begins with a letter and is followed by zero or more letters and digits

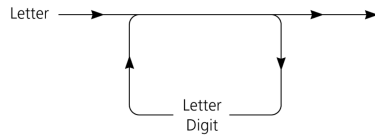


Figure 6-3

A syntax diagram for Java identifiers

© 2006 Pearson Addison-Wesley. All rights reserved

6-13

The Basics of Grammars

- Java identifiers
 - Language
 - JavaIds = {w : w is a legal Java identifier}
 - Grammar
 - < identifier > = < letter > | < identifier > < letter > | < identifier > < digit >
 - < letter > = a | b | ... | z | A | B | ... | Z | _ | \$
 - < digit > = 0 | 1 | ... | 9

© 2006 Pearson Addison-Wesley. All rights reserved

6-14

The Basics of Grammars

- Recognition algorithm

```
isId(w)
  if (w is of length 1) {
    if (w is a letter) {
      return true
    }
    else {
      return false
    }
  }
  else if (the last character of w is a letter or a digit) {
    return isId(w minus its last character)
  }
  else {
    return false
  }
```

© 2006 Pearson Addison-Wesley. All rights reserved

6-15

Two Simple Languages: Palindromes

- A string that reads the same from left to right as it does from right to left
- Examples: radar, deed
- Language
 - Palindromes = {w : w reads the same left to right as right to left}

© 2006 Pearson Addison-Wesley. All rights reserved

6-16

Palindromes

- Grammar
 - < pal > = empty string | < ch > | a < pal > a | b < pal > b | ... | Z < pal > Z
 - < ch > = a | b | ... | z | A | B | ... | Z

© 2006 Pearson Addison-Wesley. All rights reserved

6-17

Palindromes

- Recognition algorithm

```
isPal(w)
  if (w is the empty string or w is of length 1) {
    return true
  }
  else if (w's first and last characters are the same letter) {
    return isPal(w minus its first and last characters)
  }
  else {
    return false
  }
```

© 2006 Pearson Addison-Wesley. All rights reserved

6-18

Strings of the form A^nB^n

- A^nB^n
 - The string that consists of n consecutive A's followed by n consecutive B's
- Language
 $L = \{w : w \text{ is of the form } A^nB^n \text{ for some } n \geq 0\}$
- Grammar
 $\langle \text{legal-word} \rangle = \text{empty string} \mid A \langle \text{legal-word} \rangle B$

Strings of the form A^nB^n

- Recognition algorithm

```
isAnBn(w)
  if (the length of w is zero) {
    return true
  }
  else if (w begins with the character A and ends
           with the character B) {
    return isAnBn(w minus its first and last
                  characters)
  }
  else {
    return false
  }
```

Algebraic Expressions

- Three languages for algebraic expressions
 - Infix expressions
 - An operator appears between its operands
 - Example: $a + b$
 - Prefix expressions
 - An operator appears before its operands
 - Example: $+ a b$
 - Postfix expressions
 - An operator appears after its operands
 - Example: $a b +$

Algebraic Expressions

- To convert a fully parenthesized infix expression to a prefix form
 - Move each operator to the position marked by its corresponding open parenthesis
 - Remove the parentheses
 - Example
 - Infix expression: $((a + b) * c)$
 - Prefix expression: $* + a b c$

Algebraic Expressions

- To convert a fully parenthesized infix expression to a postfix form
 - Move each operator to the position marked by its corresponding closing parenthesis
 - Remove the parentheses
 - Example
 - Infix form: $((a + b) * c)$
 - Postfix form: $a b + c *$

Algebraic Expressions

- Prefix and postfix expressions
 - Never need
 - Precedence rules
 - Association rules
 - Parentheses
 - Have
 - Simple grammar expressions
 - Straightforward recognition and evaluation algorithms

Prefix Expressions

- Grammar

```
<prefix> = <identifier> | <operator> <prefix> <prefix>  
<operator> = + | - | * | /  
<identifier> = a | b | ... | z
```

- A recognition algorithm

```
isPre()  
  size = length of expression strExp  
  lastChar = endPre(0, size - 1)  
  if (lastChar >= 0 and lastChar == size-1 {  
    return true  
  }  
  else {  
    return false  
  }  
}
```

© 2006 Pearson Addison-Wesley. All rights reserved

6-25

Prefix Expressions

- An algorithm that evaluates a prefix expression

```
evaluatePrefix(strExp)  
  ch = first character of expression strExp  
  Delete first character from strExp  
  if (ch is an identifier) {  
    return value of the identifier  
  }  
  else if (ch is an operator named op) {  
    operand1 = evaluatePrefix(strExp)  
    operand2 = evaluatePrefix(strExp)  
    return operand1 op operand2  
  }  
}
```

© 2006 Pearson Addison-Wesley. All rights reserved

6-26

Postfix Expressions

- Grammar

```
<postfix> = <identifier> | <postfix> <postfix> <operator>  
<operator> = + | - | * | /  
<identifier> = a | b | ... | z
```

- At high-level, an algorithm that converts a prefix expression to postfix form

```
if (exp is a single letter) {  
  return exp  
}  
else {  
  return postfix(prefix1) + postfix(prefix2) +  
  operator  
}
```

© 2006 Pearson Addison-Wesley. All rights reserved

6-27

Postfix Expressions

- A recursive algorithm that converts a prefix expression to postfix form

```
convert(pre)  
  ch = first character of pre  
  Delete first character of pre  
  if (ch is a lowercase letter) {  
    return ch as a string  
  }  
  else {  
    postfix1 = convert(pre)  
    postfix2 = convert(pre)  
    return postfix1 + postfix2 + ch  
  }  
}
```

© 2006 Pearson Addison-Wesley. All rights reserved

6-28

Fully Parenthesized Expressions

- To avoid ambiguity, infix notation normally requires
 - Precedence rules
 - Rules for association
 - Parentheses
- Fully parenthesized expressions do not require
 - Precedence rules
 - Rules for association

© 2006 Pearson Addison-Wesley. All rights reserved

6-29

Fully Parenthesized Expressions

- Fully parenthesized expressions
 - A simple grammar
 - <infix> = <identifier> | (<infix> <operator> <infix>)
 - <operator> = + | - | * | /
 - <identifier> = a | b | ... | z
 - Inconvenient for programmers

© 2006 Pearson Addison-Wesley. All rights reserved

6-30

The Relationship Between Recursion and Mathematical Induction

- A strong relationship exists between recursion and mathematical induction
- Induction can be used to
 - Prove properties about recursive algorithms
 - Prove that a recursive algorithm performs a certain amount of work

The Correctness of the Recursive Factorial Method

- Pseudocode for a recursive method that computes the factorial of a nonnegative integer n

```
fact(n)
  if (n is 0) {
    return 1
  }
  else {
    return n * fact(n - 1)
  }
```

The Correctness of the Recursive Factorial Method

- Induction on n can prove that the method `fact` returns the values

$$\text{fact}(0) = 0! = 1$$

$$\text{fact}(n) = n! = n * (n - 1) * (n - 2) * \dots * 1 \quad \text{if } n > 0$$

The Cost of Towers of Hanoi

- Solution to the Towers of Hanoi problem

```
solveTowers(count, source, destination, spare)
  if (count is 1) {
    Move a disk directly from source to destination
  }
  else {
    solveTowers(count-1, source, spare, destination)
    solveTowers(1, source, destination, spare)
    solveTowers(count-1, spare, destination, source)
  }
```

The Cost of Towers of Hanoi

- Question
 - If you begin with N disks, how many moves does `solveTowers` make to solve the problem?
- Let
 - $\text{moves}(N)$ be the number of moves made starting with N disks
- When $N = 1$
 - $\text{moves}(1) = 1$

The Cost of Towers of Hanoi

- When $N > 1$
 - $\text{moves}(N) = \text{moves}(N - 1) + \text{moves}(1) + \text{moves}(N - 1)$
- Recurrence relation for the number of moves that `solveTowers` requires for N disks
 - $\text{moves}(1) = 1$
 - $\text{moves}(N) = 2 * \text{moves}(N - 1) + 1 \quad \text{if } N > 1$

The Cost of Towers of Hanoi

- A closed-form formula for the number of moves that `solveTowers` requires for N disks
 $\text{moves}(N) = 2^N - 1$, for all $N \geq 1$
- Induction on N can provide the proof that $\text{moves}(N) = 2^N - 1$

Summary

- Backtracking is a solution strategy that involves both recursion and a sequence of guesses that ultimately lead to a solution
- A grammar is a device for defining a language
 - A language is a set of strings of symbols
 - A recognition algorithm for a language can often be based directly on the grammar of the language
 - Grammars are frequently recursive

Summary

- Different languages of algebraic expressions have their relative advantages and disadvantages
 - Prefix expressions
 - Postfix expressions
 - Infix expressions
- A close relationship exists between mathematical induction and recursion
 - Induction can be used to prove properties about a recursive algorithm